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De Bruijn Isomorphisms and Free Space Optical Networks

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Abstract

*The de Bruijn digraph $B(d, D)$ is usually defined by words of size D on an alphabet of cardinality d , through a cyclic left shift permutation on the words, after which the rightmost symbol is changed. In this paper, we show that any digraph defined on words and alphabets of the same size, through an arbitrary permutation on the alphabet **and** an arbitrary permutation on the word indices, is isomorphic to the de Bruijn, provided that this latter permutation is cyclic. As an application, we recall that it is known that the Optical Transpose Interconnection System from UCSD can implement the de Bruijn interconnections for n nodes, for a fixed d , with $O(n)$ lenses. We show here how to improve this hardware requirement to $\Theta(\sqrt{n})$.*

Keywords: *Distributed systems, Optical networks, Free space of optical interconnections, Optical Transpose Interconnection System, de Bruijn, Graph isomorphisms, Digraphs.*

1. Introduction

Many results exist in the literature describing technological and theoretical advances in optical network topologies and design. However, an essential effort has yet to be done in linking those results together. In this paper, we propose a step in this direction, by giving optimal optical layouts of several de Bruijn-like topologies, using existing optical technologies, namely the Optical Transpose Interconnecting System (OTIS) architecture, from the Optoelectronic Computing Group of the UCSD [6, 25]. These results are obtained using graph theoretical properties that we prove in this paper. In fact, our optical layouts are particular cases of more general graph isomorphisms that we develop here.

Our motivation stems from the fact that it has been shown that the break-even line length where optical communication lines become more effective than their electrical counterparts is less than 1cm, in terms of speed and power consumption [16]. Therefore, the use of optical interconnections on-board is nowadays justified, and some studies even suggest that on-chip optical interconnects will soon be cost-effective [33]. Moreover, the emergence of cutting-edge technologies as Vertical Cavity Surface-Emitting Lasers (VCSELs) [15, 31], high sensibility optical transimpedance receivers [5], beam splitters [17, 18], micro-lenses [26], and holograms [7], makes possible the fabrication of complex optical communication networks.

The OTIS architecture is a simple scalable means of implementing very dense one-to-one interconnections between processors in a free space of optical interconnections. In a nutshell, it consists of two lenslet arrays allowing a large number of optical interconnections from a set of transmitters to a set of receivers. The OTIS architecture is used by the 3-D OptoElectronic Stacked Processors consortium to connect a 16-by-16 switch (For more details, see <http://soliton.ucsd.edu/3doesp/public/>).

It has been shown in [34] that the OTIS architecture can be used to build all-optical complete networks, based on a symmetric complete digraph with loops (K_n^*). The OTIS architecture and recent progress made in optoelectronics [24] allow the layout of a 64 processors complete digraph, each processor having 64 transceivers (corresponding to the 64 arcs of one vertex). It has also been shown in [14] how to realize the single-hop multi-OPS POPS network [10], and the multi-hop multi-OPS stack-Kautz network [13] with the OTIS architecture. However, these layouts scale badly, since they require a large number of transceivers per processor, as well as a large number of optical lenses. Consequently, an important research direction is to study the set of network

topologies for which there are efficient layouts with the OTIS architecture.

In our work, we focus on de Bruijn-like digraphs, since the de Bruijn [1] is a well known digraph which has been extensively studied in graph theory, with applications in parallel and distributed computing. A de Bruijn digraph $B(d, D)$ is an alphabet digraph of constant degree d and diameter D having for vertices the d^D words of length D on the alphabet $Z_d = \{0, 1, \dots, d-1\}$, such that a vertex $x = x_{D-1} \dots x_1 x_0$, $x_i \in Z_d$, is adjacent to the set of vertices $\Gamma^+(x) = \{x_{D-2} \dots x_1 x_0 \beta, \beta \in Z_d\}$ [1]. Many results exist on routing [19], broadcasting [28], gossiping [3] and embedding [9] on this digraph as well as on parallel algorithms for it [12, 30]. Furthermore, it has been used to build various multi-stage interconnection networks like Butterfly [30], Shufflenet [27] and GEMNET [27] and also to build fully scalable networks of any size [22] and very large bus networks [2]. More interestingly, the Galileo space probe of the NASA uses a de Bruijn network to implement a signal decoder [11].

From a theoretical point of view, we show in this paper that there are many ways to build digraphs isomorphic to the de Bruijn. For this, we start showing, in Proposition 3.2, that the alphabet digraph $B_\pi(d, D)$, built using a permutation π on Z_d and the adjacency relation $\Gamma_{B_\pi(d, D)}^+(x) = \{\pi(x_{D-2}) \dots \pi(x_1) \pi(x_0) \pi(\beta), \beta \in Z_d\}$, is isomorphic to $B(d, D)$. As a corollary, it follows the isomorphism between the de Bruijn digraph and the Imase and Itoh digraph [20], a digraph issued from the minimum-diameter maximum-size problem and which is usually said to be a generalization of the Kautz digraph [23]. Then, we show, in Proposition 3.9, that there is a permutation f on Z_D which defines an alphabet digraph, isomorphic to $B(d, D)$, if and only if f is cyclic. Notice that the permutation π is applied on the alphabet and f is applied on the word indices (see Section 3.2). Finally, we show that the permutations π and f may be applied together, which characterizes a class of digraphs which are isomorphic to the de Bruijn.

As an application of these results, we recall that it has been proved that the Imase and Itoh digraph, $II(d, n)$, with n nodes of degree d , can be implemented by the OTIS architecture [14]. Our result on the isomorphism of $II(d, n)$ and $B(d, D)$ yields that OTIS can also implement the de Bruijn. Unfortunately, for a fixed degree d , such a layout uses $O(n)$ lenses to interconnect n processors, which is not very efficient. Our subsequent results help us to show how OTIS can implement the de Bruijn with only $\Theta(\sqrt{n})$ lenses, which is not only optimal, but also very efficient for practical values, as it will be discussed later.

This paper is organized as follows. In Section 2, we give the graph theoretical machinery which we will need in the following. In particular, we remind the definition of the de Bruijn digraph and its generalization, the Reddy-Raghavan-Kuhl digraph [29]. We also recall the definition of the Imase and Itoh digraph, known to be a Kautz digraph generalization. Then, in Section 3, we show that there are many isomorphisms of the de Bruijn digraph, using first a permutation π on the alphabet Z_d and then a permutation f on the indices Z_D . We hence obtain $d!(D-1)!$ such isomorphic digraphs. The application to optical networks is given in Section 4, where we present OTIS layouts of de Bruijn-like digraphs and we show how to obtain such layouts which minimize the number of lenses used. We close the paper with some concluding remarks and directions for further research.

2. Preliminaries

2.1. Notations

As usually, Z_n will denote the ring of integers modulo n . $Z_n = \{0, 1, 2, \dots, n-1\}$. Z_n^k will denote the k dimensional vector space on Z_n and e_0, e_1, \dots, e_{n-1} the canonical base. Addition of elements of Z_n will always be performed in Z_n (i.e modulo n).

For a permutation f and an integer i , we define f^i inductively by taking f^0 as the identity permutation and $f^{i+1} = f \circ f^i$.

The next permutation will be used frequently:

Definition 2.1 *We will denote \mathcal{C} the complement permutation of Z_n defined by $\mathcal{C}(u) = n - u - 1$. As usually $\mathcal{C}(u)$ will be also denoted \bar{u} .*

2.2. Classical Digraphs

Given a digraph $G = (V, E)$, and $x \in V$ we will denote $\Gamma_G^+(x)$ the set of the out-neighbors of x , i.e. $\{y \in V, (x, y) \in E\}$. When the digraph G is clear from the context we simply denote $\Gamma^+(x)$.

Definition 2.2 [1] *The de Bruijn digraph, $B(d, D)$, of degree d and diameter D is defined as a digraph on the set of vertices made of words of length D on Z_d , and such that for $x = x_{D-1} x_{D-2} \dots x_1 x_0$, $\Gamma^+(x) = \{x_{D-2} \dots x_1 x_0 \beta, \beta \in Z_d\}$.*

Definition 2.3 *The conjunction $G_1 \otimes G_2$ of two digraphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ is the digraph with vertex-set $V_1 \times V_2$, and where there is an arc from (u_1, u_2) to (v_1, v_2) if and only if there is an both an arc from u_1 to v_1 in G_1 and an arc from u_2 to v_2 in G_2 .*

Alternatively we have $\Gamma_{G_1 \otimes G_2}^+((u_1, u_2)) = \Gamma_{G_1}^+(u_1) \times \Gamma_{G_2}^+(u_2)$.

Remark 2.4 $B(d, k) \otimes B(d', k) = B(dd', k)$; one can see [4, 32] for more details about conjunction and de Bruijn digraphs and line digraphs.

Definition 2.5 [29] The Reddy-Raghavan-Kuhl digraph, $RRK(d, n)$, of degree d with n nodes, is the digraph with vertex set Z_n and such that for $u \in Z_n$: $\Gamma^+(u) = \{du + \alpha, 0 \leq \alpha < d\}$.

Remark 2.6 The digraphs $RRK(d, d^D)$ and $B(d, D)$ are isomorphic (see [29]). To see it, we associate to a vertex $x = x_{D-1}x_{D-2} \dots x_1x_0$ of $B(d, D)$ a number $u \in Z_n$ with $u = \sum_{i=0}^{D-1} x_i d^i$; then the out-neighbors of x in $B(d, D)$ are associated to numbers of the form $du + \alpha, 0 \leq \alpha < d$. Note also that this congruent definition is sometimes considered as the standard one for the de Bruijn digraph.

Definition 2.7 [23] The Kautz digraph, $K(d, D)$, of degree d and diameter D is defined as a digraph on the set of vertices made of words of length D on the alphabet Z_{d+1} , $x = x_{D-1} \dots x_1x_0$, with $x_i \in Z_{d+1}$ and $x_i \neq x_{i+1}$. $\Gamma^+(x) = \{x_{D-2} \dots x_1x_0\beta, \beta \neq x_0, \beta \in Z_{d+1}\}$.

Definition 2.8 [20] The Imase and Itoh digraph, $II(d, n)$, of degree d with n nodes, is the digraph with vertex set Z_n and such that for $u \in Z_n$: $\Gamma^+(u) = \{-du - \alpha, 1 \leq \alpha \leq d\}$.

It has been shown in [21] that $II(d, d^{D-1}(d+1))$ is isomorphic to $K(d, D)$.

3. Alternative definition of $B(d, D)$ as a digraph on alphabet

In this section we define generalizations of de Bruijn digraph. In order to obtain the out-neighbors of a vertex labeled by some word of length D over Z_d we proceed as follows:

- instead of performing a left-shift on the word letters we permute the letter positions in some arbitrary way;
- then each letter is replaced by another one according to some fixed permutation of the alphabet Z_d ;
- finally the letter of a given fixed position is replaced by an arbitrary letter of Z_d .

3.1. Permutation on the alphabet

We start by considering only the case of a permutation on the alphabet. We will show that there is an isomorphism between $B(d, D)$ and digraphs that we call $B_\pi(d, D)$, i.e., which are similar to $B(d, D)$ up to a permutation on the alphabet Z_d . Then we will prove that $II(d, d^D)$ and $B(d, D)$ are isomorphic, which, according to the Remark 2.6 will imply that $II(d, d^D), B(d, D), RRK(d, d^D)$ are isomorphic.

Definition 3.1 Given a permutation π of Z_d , $B_\pi(d, D)$ is defined as the digraph on the set of vertices made of words of length D on Z_d , such that for each vertex $x = x_{D-1}x_{D-2} \dots x_1x_0$, $x_i \in Z_d$, $\Gamma^+(x) = \{\pi(x_{D-2}) \dots \pi(x_0)\beta, \beta \in Z_d\}$.

Proposition 3.2 $B(d, D)$ and $B_\pi(d, D)$ are isomorphic.

Proof: We just have to define an isomorphism W from $B_\pi(d, D)$ to $B(d, D)$; W is the following one to one mapping from the words of length D on the alphabet Z_d onto itself:

$$\begin{aligned} W(x_{D-1}x_{D-2} \dots x_1x_0) \\ = \pi^0(x_{D-1})\pi^1(x_{D-2}) \dots \pi^{D-2}(x_1)\pi^{D-1}(x_0) \end{aligned}$$

Given a vertex x of $B_\pi(d, D)$ we have:

$$\begin{aligned} W^{-1}\Gamma_{B(d, D)}^+(W(x_{D-1}x_{D-2} \dots x_1x_0)) \\ = W\Gamma_{B(d, D)}^+(\pi^0(x_{D-1})\pi^1(x_{D-2}) \dots \pi^{D-1}(x_0)) \\ = W^{-1}(\pi^1(x_{D-2}) \dots \pi^{D-2}x_1\pi^{D-1}(x_0)\beta), \beta \in Z_d \\ = \pi(x_{D-2}) \dots \pi(x_0)\pi(\beta) \\ = \pi(x_{D-2}) \dots \pi(x_0)\alpha, \alpha \in Z_d. \end{aligned}$$

As the last line corresponds to the adjacency relation in $B_\pi(d, D)$ we have defined the required isomorphism. \square

Notice that more generally, the digraph of degree d , diameter D and $n = d^D$ nodes, such that $\Gamma^+(x_{D-1}x_{D-2} \dots x_1x_0) = \{\pi_0(x_{D-2})\pi_1(x_{D-3}) \dots \pi_{D-2}(x_0)\pi_{D-1}(\alpha), \alpha \in Z_d\}$, where $\pi_i, i \in Z_D$, is a permutation on Z_d , is isomorphic to $B(d, D)$.

Proposition 3.3 The digraphs $B(d, D)$ and $II(d, d^D)$ are isomorphic.

Proof: Let $\overline{B}(d, D)$ denote $B_C(d, D)$, where C is the complement permutation over Z_d . As in the case of $B(d, D)$, $\overline{B}(d, D)$ can be alternatively defined as a digraph with vertex set Z_n (see Remark 2.6) in which for $u \in Z_n, \Gamma^+(u) = \{-(du + \beta) - 1, 0 \leq \beta < d\}$, that

is $\Gamma^+(u) = \{-du - \alpha, 1 \leq \alpha \leq d\}$, This is the definition of $II(d, d^D)$. So $B_C(d, D)$ and $II(d, d^D)$ are isomorphic. From proposition 3.2 these digraphs are hence isomorphic to $B(d, D)$. \square

By Remark 2.6, we obtain:

Corollary 3.4 *The digraphs $B(d, D)$, $RRK(d, d^D)$ and $II(d, d^D)$ are isomorphic.*

To illustrate our purpose, Figure 1 shows the digraph $B(2, 3)$ which is isomorphic to $RRK(2, 8)$ of Figure 2 and to $II(2, 8)$ represented on Figure 3.

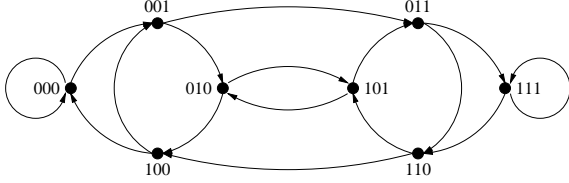


Figure 1. de Bruijn of degree 2, diameter 3 and 8 nodes.

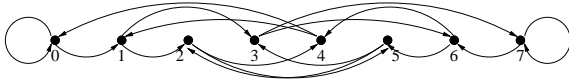


Figure 2. RRK digraph of degree 2, diameter 3 and 8 nodes.

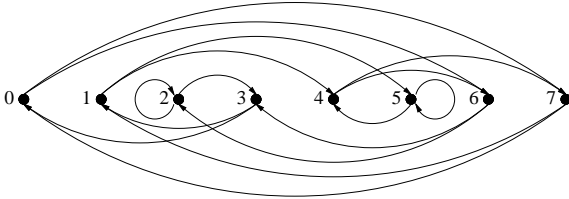


Figure 3. II digraph of degree 2, diameter 3 and 8 nodes.

3.2. Permutation on the indices, digraphs $A(f, \pi, s)$

Now we consider more general digraphs than $B_\pi(d, D)$ for which the left shift is replaced by any permutation of Z_D . First we need some definitions.

Definition 3.5 To a permutation f of Z_D , we associate the linear map \vec{f} of Z_d^D defined on the basis by $\vec{f}(e_i) = e_{f(i)}$ (this linear map is associated with a permutation matrix). Note that $\vec{f}g = \vec{f} \vec{g}$.

Definition 3.6 Any permutation π on Z_d , is extended naturally to Z_d^D by setting for $x = x_{D-1}x_{D-2} \cdots x_1x_0$: $\pi(x) = \pi(x_{D-1})\pi(x_{D-2}) \cdots \pi(x_1)\pi(x_0)$.

Definition 3.7 Let $d, D \in \mathbb{N}$, given a permutation f on Z_D , a permutation π of Z_d and a position $j \in Z_D$, we define the digraph $G = A(f, \pi, j)$ by:

$$\begin{aligned} V(G) &= Z_d^D \\ \Gamma_G^+(x) &= \pi(f(x)) + Z_d \cdot e_j \end{aligned}$$

$A(f, \pi, j)$ has degree d and we will say that it has dimension D .

Remark 3.8 Note that vertices of $B(d, D)$ can be considered as elements of Z_d^D (by associating $x_{D-1} \cdots x_1x_0$ to $\sum x_i e_i$). Let σ be the permutation of Z_D defined by $\sigma : i \mapsto i + 1 \mod D$, and Id be the identity permutation on Z_d then the de Bruijn digraph is exactly $A(\sigma, Id, 0)$. Indeed for $x = x_{D-1}x_{D-2} \cdots x_1x_0$, $\Gamma_{A(\sigma, Id, 0)}^+(x) = \{x_{D-2} \cdots x_1x_0(x_{D-1} + \gamma), \gamma \in Z_d\}$, that is also $\{x_{D-2} \cdots x_1x_0\beta, \beta \in Z_d\}$. Similarly, $B_\pi(d, D)$ and $A(\sigma, \pi, 0)$ are isomorphic.

Proposition 3.9 The graph $G = A(f, \pi, j)$ with degree d and dimension D is isomorphic to $B(d, D)$ if and only if f is a cyclic permutation of Z_D . Otherwise $A(f, \pi, j)$ is not connected.

Proof: Note that since f is a cyclic permutation it induces a unique orbit on Z_D . We associate to f the unique permutation g of Z_D , such that $\forall i \in Z_D, g(i) = f^i(j)$. From definition we have

1. $g^{-1}fg(i) = g^{-1}ff^i(j) = g^{-1}f^{i+1}(j) = g^{-1}g(i+1) = i+1$
2. $g^{-1}(j) = 0$ (Hence $\vec{g}^{-1}(e_j) = e_0$)

Let us consider the digraph isomorphism induced by \vec{g} . For $x \in Z_d^D$, we compute $\vec{g}^{-1}\Gamma_G^+(\vec{g}(x))$. From definition of G we have $\Gamma_G^+(x) = \pi(\vec{f}(x)) + Z_d \cdot e_j$, thus by linearity we obtain $\vec{g}^{-1}\Gamma_G^+(\vec{g}(x)) = \pi(\vec{g}^{-1}(\vec{f} \vec{g}(x))) + Z_d \vec{g}^{-1}(e_j)$. But from 1) $\vec{g}^{-1}fg = \vec{\sigma}$ and from 2) that $\vec{g}^{-1}(e_j) = e_0$. Finally, $\vec{g}^{-1}\Gamma_G^+(\vec{g}(x)) = \pi(\vec{\sigma}(x)) + Z_d \cdot e_0$. And this adjacency relation is the one of $B_\pi(d, D)$ which is isomorphic to $B(d, D)$. We omit the proof that $A(f, \pi, j)$ is connected only if f is cyclic. \square

Notice that since there are $(D - 1)!$ cyclic permutation functions f on Z_D , Proposition 3.9 leads to $(D - 1)!$ alternative ways of defining $B(d, D)$. Furthermore, as Proposition 3.2 leads to $d!$ other ways of defining $B(d, D)$, using permutations π on Z_d , we finally obtain $d!(D - 1)!$ alternative ways of defining $B(d, D)$.

Remark 3.10 *If f is not cyclic, $A(f, \pi, s)$ is not connected. Moreover, it can be proved that each connected component of $A(f, \pi, s)$, is the conjunction of a de Bruijn digraph with a circuit.*

3.3. Examples

For the sake of illustration of Proposition 3.9, the following examples show, first, an alphabet digraph $H = A(f, Id, 2)$ which is isomorphic to $B(d, 6)$, and then an alphabet digraph $H = A(f, Id, 1)$ where the permutation f is not cyclic and hence H is not connected.

3.3.1. $H = A(f, Id, 2)$

Let H be the digraph of degree d , dimension 6 and vertex set Z_d^6 , defined below:

$$\Gamma_H^+(x_5x_4x_3x_2x_1x_0) = x_2x_1x_0\beta x_5x_4 \text{ with } \beta \in Z_d.$$

H is indeed $A(f, Id, 2)$, where f is the following permutation of Z_6 :

$$f(i) = \begin{cases} i + 3 & \text{if } i < 3; \\ 2 & \text{if } i = 3; \\ i + 2 \pmod{6} & \text{otherwise.} \end{cases}$$

Then, \vec{f} is defined on Z_d^6 by:

$$\vec{f}(x_5x_4x_3x_2x_1x_0) = x_2x_1x_0x_3x_5x_4$$

The permutation g associated with f , used in Proposition 3.9 (i.e., $g(i) = f^i(2)$), is such that $g(0) = 2$, $g(1) = 5$, $g(2) = 1$, $g(3) = 4$, $g(4) = 0$, $g(5) = 3$, and $\vec{g}(x_5x_4x_3x_2x_1x_0) = x_1x_3x_5x_0x_2x_4$, and hence $\vec{g}^{-1}(x_5x_4x_3x_2x_1x_0) = x_3x_0x_4x_1x_5x_2$. See Figure 4 for an illustration.

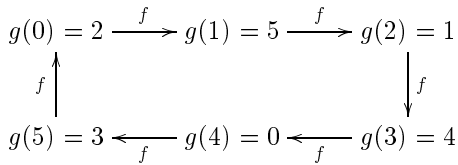


Figure 4. Illustration of the permutation g defined as $g(i) = f^i(2)$.

Then,

$$\begin{aligned} \vec{g}^{-1}\Gamma_H^+\vec{g}(x_5x_4x_3x_2x_1x_0) &= \vec{g}^{-1}\Gamma_H^+(x_1x_3x_5x_0x_2x_4) \\ &= \vec{g}^{-1}(x_0x_2x_4\beta x_1x_3), \beta \in Z_d \\ &= x_4x_3x_2x_1x_0\beta \\ &= \Gamma_{B(d,6)}^+(x_5x_4x_3x_2x_1x_0). \end{aligned}$$

Therefore, H is isomorphic to $B(d, 6)$.

3.3.2. $H = A(f, Id, 1)$

Let H be the digraph of degree d , dimension 3, and vertex set Z_d^3 , defined below:

$$\Gamma_H^+(x_2x_1x_0) = x_0\beta x_2, \text{ with } \beta \in Z_d.$$

H is indeed $A(f, Id, 1)$, where f is the permutation of Z_3 such that $f(i) = \bar{i} = 2 - i$. Then, \vec{f} is defined on Z_d^3 by $\vec{f}(x_2x_1x_0) = x_0x_1x_2$.

If we define the function g associated to f by $g(i) = f^i(1)$, we get $g(0) = g(1) = g(2) = 1$, hence g is not a cyclic permutation of Z_3 and H is not connected. Notice that the digraph H can be alternatively defined by

$$\Gamma_H^+([x_2x_0], [x_1]) = \{[x_0x_2], [\beta]\}$$

It is composed of $\frac{d^2-d}{2}$ digraphs $C_2 \otimes B(d, 1)$ plus d digraphs $C_1 \otimes B(d, 1)$, as shown in Figure 5 for $d = 2$.



Figure 5. The digraph H of example 3.3.2, when $d = 2$

4. OTIS layout for digraphs

In this section, we first recall the definition of the OTIS architecture, and the known digraphs having a layout with OTIS. Then, we study the particular case of the OTIS layout of de Bruijn digraphs, and, using the theoretical results developed above, we show how to obtain optimal OTIS layouts for the de Bruijn.

4.1. The OTIS architecture

$OTIS(p, q)$ is a free-space optical system which allows one-to-one communications from p groups of q

transmitters onto q groups of p receivers, using $p + q$ lenses. This architecture connects the transmitter (i, j) , $0 \leq i \leq p - 1$, $0 \leq j \leq q - 1$, to the receiver $(q - j - 1, p - i - 1)$. Optical interconnections in the OTIS architecture are realized with a pair of lenslet arrays [6] in a free space of optical interconnections as shown in Figure 6.

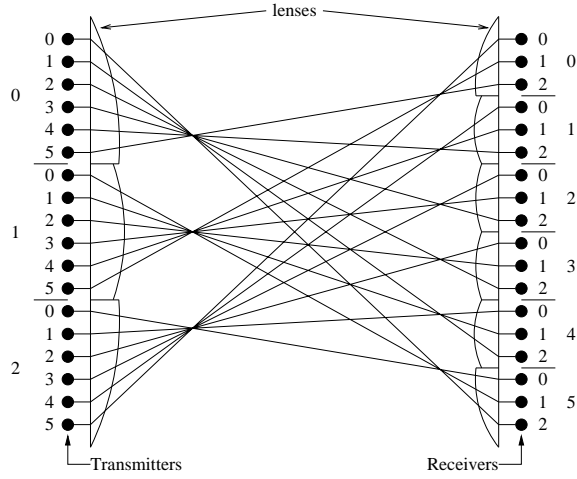


Figure 6. $OTIS(3, 6)$.

4.2. Known OTIS layouts

Let $m = pq$. The $OTIS(p, q)$ architecture connects m transmitter units of the processors to m receiver units of the same set of processors. The interconnection network can be formally defined as follows. Given a degree d such that d divides m , let $H(p, q, d)$ be the d -regular digraph with $n = \frac{m}{d}$ vertices and vertex set Z_n , in which the d transmitters $\left(\left\lfloor \frac{du + \alpha}{q} \right\rfloor, du + \alpha \bmod q \right)$, $\alpha \in Z_d$, and the d receivers $\left(\left\lfloor \frac{du + \alpha}{p} \right\rfloor, du + \alpha \bmod p \right)$, $\alpha \in Z_d$, belong to the node $u \in Z_n$ (i.e., a processing unit). There is a connection from the node u to the node v in $H(p, q, d)$ if one transmitter of u is connected to one receiver of v . For instance, figure 7 represent the digraph $H(4, 8, 2)$.

Let $G = (V, E)$ be a digraph, with $|V| = n$ nodes of constant degree d and $|E| = m = dn$ arcs. We say that G has an $OTIS(p, q)$ -layout if and only if there exist p and q with $pq = m$ and an isomorphism from $G = (V, E)$ onto $H(p, q, d)$. We remark that if G admits an $OTIS(p, q)$ -layout then G^- has an $OTIS(q, p)$ -layout, where G^- is the digraph obtained by reversing all the arcs of G .

We recall that it has been proved that the Imase and Itoh digraph, $II(d, n)$, with n nodes of degree d , has

an $OTIS(d, n)$ -layout [14], and consequently, so do the de Bruijn and the Kautz by Proposition 3.3. For fixed degree d , such a layout uses $p + q = O(n)$ lenses to interconnect n processors, which is not very efficient. In the following we will show how to build an $OTIS(p, q)$ -layout for the de Bruijn, such that the number of lenses used is $p + q = \Theta(\sqrt{n})$, which is also very efficient for practical values. We remark that the values of p and q determine also the size of the lenses and, by technological considerations, it is preferable to work with lenses of almost same size (i.e., $p \sim q$).

4.3. The degree-diameter problem for OTIS layouts

For given degree d and diameter D , one can study the maximum number of nodes of the family of digraphs $H(p, q, d)$ (i.e., the digraph of degree d built from an $OTIS(p, q)$). We obtained, by exhaustive search, the results reported in Table 1 for degree 2 and diameters 8, 9, and 10. The table contains only the largest digraphs found for each diameter.

The Kautz digraph appears to be the largest digraph of degree d and diameter D which has an $OTIS(p, q)$ -layout. This was predictable because it is among the best known constructions for the classic degree-diameter problem [8].

For fixed degree d , diameter D , and number of nodes $n = d^D$, we found several digraphs $H(p, q, d)$, with $pq = d^{D+1}$ for different values of p and q such that both p and q were powers of d . All these digraphs happened to be isomorphic to $B(d, D)$. As an example, $H(2, 256, 2)$, $H(4, 128, 2)$ and $H(16, 32, 2)$ are isomorphic to $B(2, 8)$. In the next section, we will prove that whenever D is even and $H(p, q, d)$, as defined above, is connected, then it is isomorphic to $B(d, D)$. We will also address the case where D is odd.

4.4. Efficient $OTIS(p, q)$ -layout of de Bruijn digraphs

In this section we characterize $OTIS(p, q)$ -layouts of de Bruijn digraph having a minimal number of lenses. For this, we will study cases where $p = d^{p'}$ and $q = d^{q'}$, as indicated in the discussion above.

Proposition 4.1 *The digraph $H(p, q, d)$ is isomorphic to $A(f, \mathcal{C}, p' - 1)$, where \mathcal{C} is the complement permutation as defined in Definition 2.1, and where the permutation f is defined on Z_D , $D = p' + q' - 1$, as follows:*

$$f : i \longmapsto \begin{cases} i + p' & \text{if } i < q' - 1; \\ p' - 1 & \text{if } i = q' - 1; \\ i + p' - 1 \bmod D & \text{otherwise.} \end{cases}$$

n	p	q		n	p	q	
\vdots	\vdots	\vdots		\vdots	\vdots	\vdots	
253	2	253		509	2	509	
254	2	254		510	2	510	
255	2	255		511	2	511	
256	2	256	B(2,8)	512	2	512	B(2,9)
	4	128			8	128	
	16	32		513	2	513	
258	2	258		516	2	516	
264	2	264		528	2	528	
288	2	288		576	2	576	
384	2	384	K(2,8)	768	2	768	K(2,9)
$D = 8$				$D = 9$			

n	p	q	
\vdots	\vdots	\vdots	
1022	2	1022	
1023	2	1023	
1024	2	1024	B(2,10)
	4	512	
	8	256	
	16	128	
	32	64	
1026	2	1026	
1032	2	1032	
1056	2	1056	
1152	2	1152	
1536	2	1536	K(2,10)
$D = 10$			

Table 1. $H(p, q, 2)$ with diameter 8, 9 and 10.

Proof: Notice that in this proof, we associate to a number $x \in Z_{d^k}$, its d -ary representation, denoted $\omega(x)$ (i.e., $\omega(x) = x_{k-1} \dots x_1 x_0$, $x_i \in Z_d$, such that $x = \sum_{i=0}^{k-1} x_i d^i$).

Recall that the $OTIS(p, q)$ architecture connects p groups of q transmitters to q groups of p receivers, such that the transmitter (i, j) , $0 \leq i < p$, $0 \leq j < q$, is connected to the receiver $(q - j - 1, p - i - 1)$. Remark that $\omega(i)$ has length p' and $\omega(j)$ has length q' . Then, we label transceiver (i, j) by $\omega(i)\omega(j)$, of length $p' + q'$. As $\omega(q - j - 1) = \overline{\omega(j)}$ and $\omega(p - i - 1) = \overline{\omega(i)}$, we have that the transmitter $\omega(i)\omega(j)$ is connected to the receiver $\omega(j)\omega(i)$.

$H(d^{p'}, d^{q'}, d)$ has constant degree d and n nodes, and note that since d divides both p and q , each group of $d^{q'}$ transmitters contains $d^{q'-1}$ nodes. So we can label the k -th node ($k \in Z_{d^{q'-1}}$) of the i -th group of transmitters ($i \in Z_p$) by $\omega(i)\omega(k)$, of length $p' + q' - 1$. This node has the transmitters $\{(i, dk + \beta), \beta \in Z_d\}$, which are labeled by the set of strings $\{\omega(i)\omega(k)\omega(\beta), \beta \in Z_d\}$. Thus, the transmitters of the node $\omega(i)\omega(k)$ is connected to the receivers $\{\omega(k)\omega(\beta)\omega(i), \beta \in Z_d\}$. Let $\omega(i) = \omega(l)\omega(\alpha)$, where $\omega(l)$ is a string of length $p' - 1$ and $\alpha \in Z_d$. We have $\{\omega(k)\omega(\beta)\omega(i), \beta \in Z_d\} = \{\omega(k)\omega(\beta)\omega(l)\omega(\alpha), \beta \in Z_d\}$ and these receivers belong to the nodes $\{\omega(k)\omega(\beta)\omega(l), \beta \in Z_d\}$.

Finally, a node x represented by the string $\omega(x) = \omega(l)\omega(\alpha)\omega(k)$, with $\omega(l) \in Z_d^{p'-1}$, $\omega(\alpha) \in Z_d$ and $\omega(k) \in Z_d^{q'-1}$, is connected through the $OTIS(p, q)$ architecture, to the nodes represented by the strings $\{\overline{\omega(k)\omega(\beta)\omega(l)}, \beta \in Z_d\}$. As $\omega(x) = x_{D-1} \dots x_1 x_0$, $x_i \in Z_d$, we have $\omega(l) = x_{D-1} x_{D-2} \dots x_{q'}$, $\omega(\alpha) = x_{q'-1}$, and $\omega(k) = x_{q'-2} \dots x_1 x_0$. Thus, the digraph $H(d^{p'}, d^{q'}, d)$ can be defined on the vertex set Z_d^D with

$$\Gamma_H^+(x_{D-1} \dots x_1 x_0) = \{\overline{x_{q'-2} \dots x_0 \beta x_{D-1} \dots x_{q'}}\}.$$

According to the definition of f this is exactly $A(f, \mathcal{C}, p' - 1)$. □

Corollary 4.2 Let $p' + q' - 1 = D$. For any degree d , the digraphs $B(d, D)$ and $H(d^{p'}, d^{q'}, d)$ are isomorphic if and only if the permutation f of Z_D defined by

$$f : i \mapsto \begin{cases} i + p' & \text{if } i < q' - 1; \\ p' - 1 & \text{if } i = q' - 1; \\ i + p' - 1 \pmod{D} & \text{otherwise,} \end{cases}$$

is cyclic.

Proof: We showed in Proposition 4.1 that $H(d^{p'}, d^{q'}, d)$ is isomorphic to the alphabet digraph $A(f, \mathcal{C}, p' - 1)$. We have also showed in Proposition

3.9 that the alphabet digraph $A(f, \mathcal{C}, s)$ of degree d and dimension D is isomorphic to the de Bruijn digraph $B(d, D)$ when f is cyclic. The corollary follows. \square

Now, we show that there exist p' and q' such that $B(d, D)$ and $H(d^{p'}, d^{q'}, d)$ are isomorphic and $d^{p'} + d^{q'} = \Theta(\sqrt{d^D})$ (i.e., $p + q = \Theta(\sqrt{n})$).

Proposition 4.3 *Let D be odd and $p' = q' = \frac{D+1}{2}$. Then $B(d, D)$ and $H(d^{\frac{D+1}{2}}, d^{\frac{D+1}{2}}, d)$ are isomorphic if and only if $D = 1$.*

Proof: If $p' = q' > 1$, $f(p' - 1) = q' - 1 = p' - 1$ and f is not cyclic. If $p' = q' = 1$, $D = 1$ and $B(d, 1)$ is isomorphic to $H(d, d, d)$. \square

Corollary 4.4 *Let $p' = \frac{D}{2}$ and $q' = \frac{D}{2} + 1$. For fixed a degree d , if the diameter D is even, the digraphs $B(d, D)$ and $H(d^{\frac{D}{2}}, d^{\frac{D}{2}+1}, d)$ are isomorphic, and therefore, $B(d, D)$ has an $OTIS(p, q)$ -layout with $p + q = \Theta(\sqrt{n})$ lenses.*

Proof: The permutation f on Z_D is defined by

$$f : i \mapsto \begin{cases} i + \frac{D}{2} & \text{if } i < \frac{D}{2}; \\ \frac{D}{2} - 1 & \text{if } i = \frac{D}{2}; \\ i + \frac{D}{2} - 1 \pmod{D} & \text{otherwise,} \end{cases}$$

and one can easily verify that f is cyclic. Thus, using Corollary 4.2, the isomorphism is verified.

Figure 7 represents the digraph $H(4, 8, 2)$ with vertex set Z_2^4 and such that $\Gamma_H^+(x_3 x_2 x_1 x_0) = \{x_1 x_0 \beta x_3\}$, and Figure 8 represents the digraph $B(2, 4)$ with the nodes labeling of $H(4, 8, 2)$.

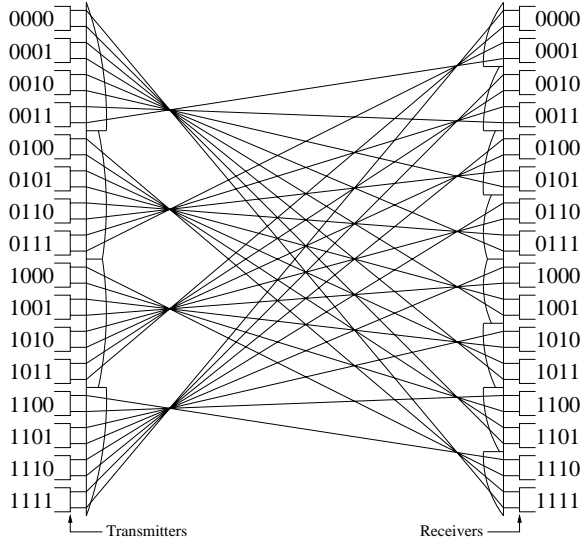


Figure 7. $H(4, 8, 2)$.

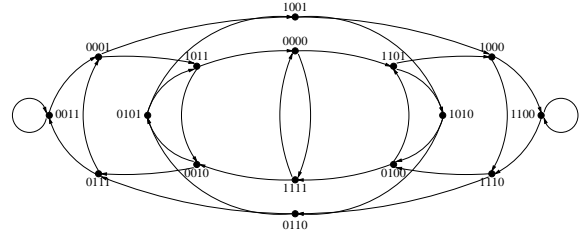


Figure 8. $B(2, 4)$ with the adjacency relation of $H(4, 8, 2)$, i.e.: $\Gamma_H^+(x_3 x_2 x_1 x_0) = \{x_1 x_0 \beta x_3\}$.

When the diameter $D > 1$ is odd, by Proposition 4.3, $B(d, D)$ and $H(p, q, d)$ cannot be isomorphic with $p = q$. The case where p is closest to q is hence when, w.l.o.g., $p' = \frac{D-1}{2}$ and $q' = \frac{D-1}{2} + 2$. However, it is not always the case that $B(d, D)$ and $H(d^{\frac{D-1}{2}}, d^{\frac{D-1}{2}+2}, d)$ are isomorphic. For instance, $H(2^5, 2^7, 2)$ and $B(2, 11)$ are isomorphic, while $H(d^6, d^8, d)$ and $B(d, 13)$ are not.

Corollary 4.5 *Given $B(d, D)$ and $H(p, q, d)$, their isomorphism can be verified in time $O(D)$.*

Proof: Using Corollary 4.2 it is sufficient to verify that f is a cyclic permutation of Z_D . This requires $O(D)$ steps. \square

Corollary 4.6 *Minimizing $d^{p'} + d^{q'}$, such that $B(d, D)$ and $H(d^{p'}, d^{q'}, d)$ are isomorphic, requires $O(D^2)$ time steps.*

Proof: If D is even, using Corollary 4.4, the minimum is obtained with $p' = \frac{D}{2}$ and $q' = \frac{D}{2} + 1$. Otherwise, there are D pairs (p', q') , with $p' < q'$, such that $D = p' + q' - 1$. Thus, using Corollary 4.5, we obtain the minimum value in time $O(D^2)$. \square

5. Conclusion

In this paper, using graph theoretical properties that we proved, we characterized the $OTIS(p, q)$ -layout of the de Bruijn digraph $B(d, D)$ when p and q are powers of d ($p = d^{p'}$, $q = d^{q'}$ and $p' + q' - 1 = D$). Consequently, in the case where d is prime, this implies a characterization of the $OTIS(p, q)$ -layout of the de Bruijn. For the other cases, one could consider $OTIS(p, q)$ -layouts with $p, q \neq d^i$, but intuition and exhaustive search make us conjecture that, except for trivial cases, such layouts do not exist.

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